

Optimizing Energy Efficiency in Bulk Transfer Networks

Matthias Herlich, Holger Karl

Problem Statement

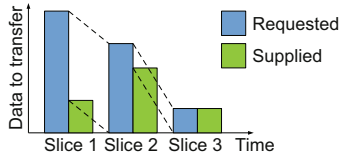
Given a network and a set of bulk transfers (instead of flows), what is the most energy-efficient route for the bulk transfers?

Background

Comparing Bulk Transfers to Flows

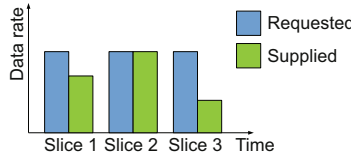
Bulk transfers

- Model for large amounts of data (e.g. CDN)
- Specify amount of data to transmit
- Done after predefined amount of data has been transmitted



Flows

- Model for streams or aggregation of many small transfers
- Specify required data rate
- Continue for an unknown time



Slices of time

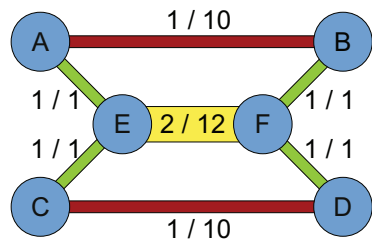
- Needed to optimize energy efficiency using mixed integer programming
- Constant data rate in a single slice
- A bulk transfer can take several slices to finish

Assumption

Network hardware can be disabled to conserve energy

Model – Definition of the Network and Bulk Transfers

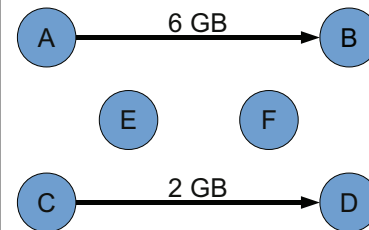
Network Graph



Undirected graph $G=(V, E)$
Capacity of edges $c: E \rightarrow \mathbb{Q}$
Energy consumption of edges $\varepsilon: E \rightarrow \mathbb{Q}$

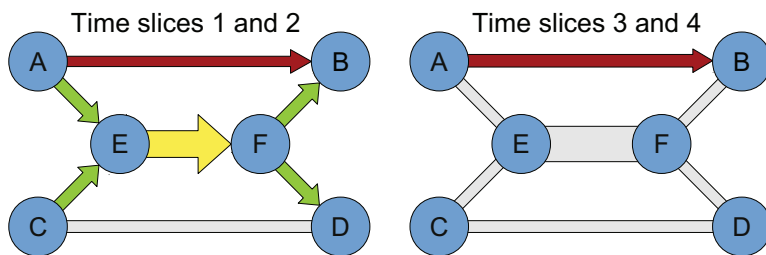
Bandwidth (width) [GB/slice] / Energy Consumption (color) [J/slice]

Requested Bulk Transfers



Set of bulk transfers: B
 $\forall b \in B: \begin{cases} \text{Source } s_b \in V \\ \text{Destination } d_b \in V \\ \text{Transfer size } a_b \in \mathbb{Q} \end{cases}$

Solution – Finding Routes with Optimal Energy Efficiency using MIP



Optimal solution needs multi-path routing!

Minimize $\sum_{t \in T} \sum_{e \in E} x^t(e) \varepsilon(e)$ (total energy consumption)

Where $x^t(e) = \begin{cases} 0, & \text{if } \forall b \in B: f_b^t(e) = 0 \\ 1, & \text{else} \end{cases}$ (edge activity)
 T set of needed slices
 $f_b^t: E \rightarrow \mathbb{Q}$ (flow transmitted for $b \in B, t \in T$)

With flow constraints:
 $\forall t \in T, e \in E: \sum_{b \in B} \max(0, f_b^t(e)) \leq c(e)$ (capacity constraint, full duplex)
 $\forall b \in B, t \in T, (u, v) \in E: f_b^t((u, v)) = -f_b^t((v, u))$ (skew symmetry)
 $\forall b \in B, t \in T, u \in V \setminus \{s_b, d_b\}: \sum_{v \in N(u)} f_b^t((u, v)) = 0$ (flow conservation)
 $\forall b \in B: \sum_{t \in T} \sum_{v \in N(s_b)} f_b^t((s_b, v)) = a_b$ (complete transfer)
(with $N(v)$ being the set of vertices incident to v)

Contact person

Matthias Herlich
+49 5251 60-5382
herlich@upb.de

University of Paderborn
Pohlweg 47-49
33098 PADERBORN
GERMANY
<http://www.cs.upb.de/cs/cn/>

Head of group

Prof. Dr. Holger Karl
+49 5251 60-5375
holger.karl@upb.de