

Prony's method in several variables, oligonomials and Linear Algebra

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Introduction

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ESSAI EXPERIMENTAL

ET ANALYTIQUE

*Sur les lois de la Dilatabilité des fluides élastiques et sur celles
de la Force expansive de la vapeur de l'eau et de la vapeur
de l'alkool, à différentes températures.*

Par R. PRONY.

Prony's problem

For a *exponential polynomial*

$$f(x) = \sum_{\omega \in \Omega} f_{\omega} e^{\omega^T x}, \quad \Omega \subset (\mathbb{R} + i\mathbb{R}/2\pi\mathbb{Z})^s, \quad 0 \neq f_{\omega} \in \mathbb{C},$$

determine Ω and f_{ω} from samples of f .

Oligonomials

For an *algebraic polynomial*

$$f(x) = \sum_{\alpha \in A} f_{\alpha} x^{\alpha}, \quad A \subset \mathbb{N}_0^s, \quad 0 \neq f_{\alpha} \in \mathbb{C}, \quad \#A < \infty,$$

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For a **sparse** *exponential polynomial*

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Oligonomials (fewnomials)

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determine A and f_{α} from samples of f .

Assumptions

- 1 Problem is *sparse*: $\#\Omega/\#A$ small.
- 2 Frequencies $\omega \in \Omega$ and powers $\alpha \in A$ can vary.
- 3 No embedding in $\Pi_{\deg A}$.
- 4 A priori knowledge: $\#\Omega$ or $\#A$.

Problem structure (Prony)

- Frequencies: *nonlinear* problem.
- Coefficients: *linear* problem.
- Evaluation points: Grids $\Gamma \subset \mathbb{Z}^s$.
- Real only – complex straightforward.

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- 1 $\Pi = \mathbb{R}[x] = \mathbb{R}[x_1, \dots, x_s]$.
- 2 $\Pi_n = \left\{ p(x) = \sum_{|\alpha| \leq n} p_\alpha x^\alpha : p_\alpha \in \mathbb{R} \right\}$ of degree n .
- 3 *Degree*

$$\deg p = \max \{ |\alpha| : f_\alpha \neq 0 \},$$

Coefficient vectors

$$p \simeq \mathbf{p} = [p_\alpha : |\alpha| \leq \deg p].$$

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A Hankel matrix

$$\mathbf{F}_n := \left[f(\alpha + \beta) : \begin{array}{l} |\alpha| \leq n \\ |\beta| \leq n \end{array} \right] \in \mathbb{R}^{d_n \times d_n}, \quad d_n = \binom{n+s}{s}.$$

A computation ...

For $p \in \Pi_n$:

Consequence

$p(X_\Omega) = 0$ implies $\mathbf{F}_n p = 0$.

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The ideal

- 1 $I_\Omega = I(X_\Omega) = \{p \in \Pi : p(X_\Omega) = 0\}.$
- 2 Zero dimensional ideal.
- 3 ...

Prony again

- 1 $p \in I_\Omega \cap \Pi_n \Rightarrow p \in \ker F_n.$
- 2 Converse?
- 3 $\Omega = \{\omega, \omega'\}, f_\omega = -f_{\omega'}$

Theorem

For n large enough we have $I_\Omega \cap \Pi_n \simeq \ker F_n.$

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- 2 Converse? No!
- 3 $\Omega = \{\omega, \omega'\}, f_\omega = -f_{\omega'}$ gives $F_0 = 0.$

Theorem

For n large enough we have $I_\Omega \cap \Pi_n \simeq \ker F_n.$

Goal

Answer the question *What is “large enough”?*

Interpolation space

$\mathcal{P} \subset \Pi$ *degree reducing interpolation space* with respect to a *finite* $X \subset \mathbb{R}^s$:
for $g \in \Pi$ there exists a *unique* $p \in \mathcal{P}$ such that

- $p(X) = g(X)$.
- $\deg p \leq \deg g$.

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If $\mathcal{P}, \mathcal{Q}$ are degree reducing interpolation spaces for X then

$$\deg \mathcal{P} = \max\{\deg p : p \in \mathcal{P}\} = \deg \mathcal{Q}$$

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Answer the question *What is “large enough”?* $n \geq d(X_\Omega)$.

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H-bases

- ① $H \subset \Pi$ *H-basis* for I if

$$g \in I \quad \Leftrightarrow \quad g = \sum_{h \in H} g_h h, \quad \deg g_h + \deg h \leq \deg g.$$

- ② Alternative: *homogeneous ideals*. (Macaulay 1916)

Definition

(Homogeneous) *leading form*

$$\Lambda(p) = p_{\deg p}^0 = \sum_{|\alpha| = \deg p} p_{\alpha} (\cdot)^{\alpha},$$

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(Homogeneous) leading form and ideal generated by G

$$\Lambda(p) = p_{\deg p}^0 = \sum_{|\alpha|=\deg p} p_\alpha (\cdot)^\alpha, \quad \langle G \rangle = \left\{ \sum_{g \in G} q_g g : q_g \in \Pi \right\}.$$

Further ingredient

- 1 Inner product $(\cdot, \cdot) : \Pi \times \Pi \rightarrow \mathbb{R}$.
- 2 For example $(p, q) = p^T q = \sum_{\alpha} p_{\alpha} q_{\alpha}$.

Division with remainder for p, G

- 1 While $p \neq 0$
 - 2 Decompose $\Lambda(p) = \sum_{g \in G} p_g \Lambda(g) + r$
 - 3 Replace $p \leftarrow p - \sum_{g \in G} p_g \Lambda(g)$
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Theorem

If H is an H -basis and $p = \sum_{g \in G} p_h h + r$ computed by reduction, then

- 1 $r = 0$ iff $p \in \langle H \rangle$.
- 2 r depends only on $\langle H \rangle$

Definition

- 1 $\nu_I(p) := r = p - \sum_{g \in G} p_h h$ normal form of p modulo $\langle H \rangle =: I$.
- 2 $N_I := \nu_I(\Pi) \simeq \Pi / I$ inverse system for I . (Macaulay '16, Gröbner '37)

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Theorem (reduction is interpolation)

$N_{I(X)}$ is a degree reducing interpolation space for X .

Normal forms for points

- 1 Given $X \subset \mathbb{R}^s$, $\#X < \infty$.
- 2 Basis for normal form space
- 3 Satisfies $\ell_x(x') = \delta_{xx'}$, $x, x' \in X$.

Theorem

For $p \in \Pi$ are equivalent:

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Goal

- 1 Compute H-basis and normal form space N_Ω for $I_\Omega = I(e^\Omega)$.
- 2 Compute frequencies Ω (standard).
- 3 Compute coefficients f_ω (linear system).

Slightly different Hankel matrix

$$F_{n,k} = \left[f(\alpha + \beta) : \begin{array}{l} |\alpha| \leq n \\ |\beta| \leq k \end{array} \right], \quad k \leq n.$$

The “magic” number

- Fundamental guess: $n > d(X_\Omega) = \deg N_\Omega$, e.g., $n \geq \#\Omega$.
- Also needed in 1d Prony!

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- Fundamental guess: $n > d(X_\Omega) = \deg N_\Omega$, e.g., $n \geq \#\Omega$.
- Also needed in 1d Prony!

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- 1 Compute H-basis and normal form space N_Ω for $I_\Omega = I(e^\Omega)$.
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Kernel vs. rank

With $F_{n,k} = U_k \Sigma V_k^T$:

- 1 $\ker F_{n,k}$: V_* columns of V for $\sigma_{\ell,\ell} = 0$.
- 2 $\text{rank } F_{n,k}$: V_+ columns wrto nonzero singular values.

The H-basis update

- 1 With

$$H_{\leq k-1} = [H_0 \dots H_{k-1}], \quad H_j = [h_{j1} \dots h_{jv_l}],$$

- 2 decompose

$$V_*^+ H_{\leq k-1} = QR = [Q_1 \ Q_2] \begin{bmatrix} R \\ 0 \end{bmatrix}.$$

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- 1 General idea: $N_k = \Pi_k^0 \ominus \Lambda(H_k)$.
- 2 Again

$$\Lambda(H_k) = Q_k \begin{bmatrix} R_k \\ 0 \end{bmatrix}, \quad Q_k = [Q_{k,1} \ Q_{k,2}].$$

- 3 $N_j = Q_{j,2}$.

Reduction

- With $k := \deg p$ solve $R_k c = Q_{k,1}^T \Lambda(p)$.
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Theorem

If $k > \deg N_\Omega$

- 1 The columns of $H_{\leq k}$ are coefficient vectors of an H-basis of I_Ω .
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Remarks

- All by means of standard Linear Algebra (Matlab).
- H-basis is *very* redundant.
- Reduction is simple and fast.
- Numerically stable due to *orthogonality*.

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If $k > \deg N_\Omega$

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The *eigenvalues* of the M_j are $(z_\omega)_j$ and the *eigenvectors* ℓ_ω , $\omega \in \Omega$.

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$$L_{k,j} N_k, \quad L_{k,j} = \sum_{|\alpha|=k} e_{\alpha+\epsilon_j} e_{\alpha}^T.$$

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2 Computation of multiplication and reduction

$$\left(I - Q_{k,1} Q_{k,1}^T \right) L_{k,j} N_k, \quad L_{k,j} = \sum_{|\alpha|=k} e_{\alpha + \epsilon_j} e_{\alpha}^T.$$

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Solve ...

$$\left[e^{\omega^T \alpha} : \begin{array}{l} |\alpha| \leq k \\ \omega \in \Omega \end{array} \right] [f_\omega : \omega \in \Omega] = [f(\alpha) : |\alpha| \leq k].$$

Properties

- Transpose of *Vandermonde matrix* for X_Ω .
- Full rank.
- Overdetermined system, stabilizes.
- Right hand side: part of $F_{n,0}$.

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- 1 Get a good guess for n .
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Random frequencies & coefficients, real, 100 tests

parameters			average error		max error	
s	# freq.	n	coeff	freq	coeff	freq
2	5	3	1.3688e-11	1.8332e-09	3.5131e-09	2.4165e-07
2	10	5	4.9366e-08	2.6388e-06	7.3010e-05	5.3330e-04
2	15	8	7.0614e-07	2.9725e-04	1.4659e-04	4.4493e-02
2	20	9	Inf	Inf	NaN	NaN
3	20	6	1.5874e-08	1.4165e-06	4.7337e-05	8.9382e-04
4	20	5	8.4712e-12	4.6565e-11	9.0309e-09	3.7456e-09
5	20	5	1.6879e-12	5.9416e-11	1.9510e-09	1.3243e-08
5	50	5	1.1079e-10	6.6070e-10	3.1709e-07	6.6913e-08
5	100	6	2.9307e-09	1.9431e-08	1.0034e-05	1.3912e-06
5	150	8	1.3142e-08	8.4199e-08	5.7281e-06	4.3975e-06

Random frequencies, purely imaginary, 100 tests

parameters			average error		max error	
s	# freq.	n	coeff	freq	coeff	freq
2	10	5	1.3476e-14	3.4744e-13	6.0290e-12	1.3724e-10
2	20	7	2.5148e-14	1.2420e-12	3.2103e-11	7.8847e-10
2	50	11	5.9357e-14	3.9721e-12	1.1845e-10	5.5214e-09
2	100	15	9.0480e-13	5.7684e-11	8.8308e-09	2.0468e-07
5	100	6	2.3796e-15	4.3794e-15	3.1431e-11	3.2918e-14
5	150	8	2.3954e-15	4.7773e-15	1.1702e-11	6.9726e-14

Line through origin, top real, bottom imaginary, 100 tests

parameters			average error		max error	
s	# freq.	fail	coeff	freq	coeff	freq
2	3	0	5.1668e-06	1.6241e-04	0.0023195	0.0243550
2	4	0	2.8912e-06	2.9318e-03	9.9505e-04	5.8547e-01
2	5	5	3.1901e-05	2.1641e-02	0.0058405	1.8753920
2	10	100	∅	∅	∅	∅
3	5	5	1.4484e-05	2.8744e-02	0.0016677	1.5547492
3	4	14	2.1197e-05	1.2439e-01	4.2678e-03	1.8590e+01
3	5	24	3.8699e-04	3.9617e-02	0.057250	1.326782
2	5	0	2.1330e-12	3.6481e-11	3.9225e-10	4.5345e-09
2	10	0	3.1867e-06	5.9222e-03	0.0018326	1.7269961
2	20	11	8.0145e-06	4.2270e-03	0.0071972	1.0316399
3	10	1	0.0025437	0.0206981	1.0404	4.3874

Recall

Sparse polynomials

$$f(x) = \sum_{\alpha \in A} f_{\alpha} x^{\alpha}.$$

Oligo $\simeq$ “few”

Simple transformation

- Pick any invertible $\Xi \in \mathbb{C}^{s \times s}$.
- Consider

$$F_{n,k}(\Xi) := \left[e^{\Xi(\beta+\gamma)} : \begin{array}{l} |\beta| \leq n \\ |\gamma| \leq k \end{array} \right].$$

- Prony situation with $\Omega = \Xi^T A = \{\Xi^T \alpha : \alpha \in A\}$:
- Round to next integer: $A = \lceil \Xi^{-T} \Omega \rceil$.

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- 3 Prony situation with $\Omega = \Xi^T A = \{\Xi^T \alpha : \alpha \in A\}$: *exponential grid*.
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$$F_{n,k}(\Xi) := \left[e^{\Xi(\beta+\gamma)} : \begin{array}{l} |\beta| \leq n \\ |\gamma| \leq k \end{array} \right].$$

- 3 Prony situation with $\Omega = \Xi^T A = \{\Xi^T \alpha : \alpha \in A\}$: *exponential grid*.
- 4 Round to next integer: $A = \lceil \Xi^{-T} \Omega \rceil$.

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Thank you!